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# Why Statistical Agencies Need to Take Privacy-loss Budgets Seriously, and What It Means When They Do

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# Why Statistical Agencies Need to Take Privacy-loss Budgets Seriously, and What It Means When They Do

## **Comments**

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At the time this talk was given, Abowd was Associate Director for Research and Methodology and Chief Scientist, U.S. Census Bureau. The opinions expressed in this talk are his own.

# Why Statistical Agencies Need to Take Privacy-loss Budgets Seriously, and What It Means When They Do

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2016 FCSM Statistical Policy Seminar  
*The Future of Federal Statistics – Use of Multiple Data Sources, Anchored  
in Fundamental Principles and Practices*  
December 6-7, 2016

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- The opinions expressed in this talk are the my own

# Outline

- The database reconstruction theorem, a.k.a. the fundamental law of information recovery
- What is a privacy-loss budget?
- How do you respect a privacy-loss budget?
- How do you prove that the rate of privacy loss in published data is consistent with the budget?
- What does it mean to prove that the released data are robust to all future attacks?

# The Database Reconstruction Theorem

- Powerful result from Dinur and Nissim (2003) [[link](#)]
- *Too many statistics published too accurately from a confidential database exposes the entire database with certainty*
- How accurately is “too accurately”?
  - Cumulative noise must be of the order  $\sqrt{N}$

# Database Reconstruction II

- Led quickly to “differential privacy”:
  - Dwork, McSherry, Nissim, and Smith (2006) [[link](#)]
  - Dwork (2006) [[link](#)]
- Leading formal privacy model

# Database Reconstruction III

- “The Fundamental Law of Information Recovery”
  - Dwork and Roth, 2014 [[link](#)]
  - Dwork, undated [[link](#)]
- Includes extensions found in
  - Dwork, McSherry and Talwar (2007) [[link](#)]
  - Muthukrishnan and Nikolov (2012) [[link](#)]
  - Kasiviswanathan, Rudelson and Smith (2013) [[link](#)]
  - Dwork, Smith, Steinke, Ullman, and Vadhan (2015) [[link](#)]



# Historical Note

- The U.S. Census Bureau: first organization in the world to use a formally private confidentiality protection system in production
  - [OnTheMap](#) (residential side)
- Machanavajjhala, Kifer, Abowd, Gehrke, and Vilhuber (2008)  
[\[link\]](#)

# What is a Privacy-loss Budget?

- Not a dollar budget, but works the same way
- Constrains aggregate risk of partial database reconstruction given all published statistics
- Worst-case limit to the inferential disclosure of any identity or item
- In differential privacy, worst case is over all possible databases with the same schema for all individuals and items

# Why Use Worst-case Protection?

- “Worst case” is “equal protection under the law”
  - Protects every person in the population the same way
  - Anyone who might have been selected for the census or survey, whether in the database or not
- “Average-case” protection does not
  - Can identify who is advantaged or disadvantaged *a priori*

# Respecting a Privacy-loss Budget

- All released statistics can *never* permit a database reconstruction more accurate than the budget
- Protection into the indefinite future
- For differential privacy, guarantee is over all future attackers and any database with the same schema

# Current Context

- Don't current confidentiality laws require data stewards to respect a privacy-loss budget, at least implicitly?
- Unclear
- Law are silent on limitations of what can be learned about the confidential data from the released statistics (database reconstruction)
- All data publication inherently involves some inferential disclosure risk; otherwise, it is useless
  - Dwork and Naor (2008) [[link](#)]: impossibility theorem
  - Kifer and Machanavajjhala (2011) [[link](#)]: no free lunch theorem

# This Is Not a New Problem

- Ratio of the circumference of a circle to its diameter is constant
- Ancients didn't understand irrational numbers:
  - Babylonians:  $\pi = 3 \frac{1}{8}$
  - Egyptians:  $\pi = 4 \times (\frac{8}{9})^2$
  - Israelites:  $\pi = 3$  [Talmud legislated value]
  - Hindu:  $\pi = \frac{62,832}{20,000} = 3.1416$
  - Euclid: no rational number is exact for this problem
  - Archimedes: sequences can approximate  $\pi$  with increasing accuracy
- But legal documents continued to use crude approximations
- Takes time to process abstract ideas into practical laws
- Legal guidance on inferential disclosure limitation is important
- But must be constructed sensibly

Source: Beckman, Petr "A History of Pi" (1971) [\[link\]](#)

# Example: Randomized Response

- Randomized response is provably privacy-loss protective
- Privacy loss bounded by the maximum Bayes factor

$$\max BF = \frac{\frac{Pr[SQ = Yes|A = Yes]}{Pr[SQ = No|A = Yes]}}{\frac{Pr[SQ = Yes]}{Pr[SQ = No]}} = \frac{Pr[A = Yes|SQ = Yes]}{Pr[A = Yes|SQ = No]} = \frac{(1/2) + (1 - 1/2)^{1/2}}{(1 - 1/2)^{1/2}} = 3$$

- Bound is the logarithm of the maximum Bayes factor
- If
  - Sensitive question asked with probability  $\frac{1}{2}$
  - And innocuous question is “yes” with probability  $\frac{1}{2}$
  - Then the maximum Bayes factor is 3, and  $\ln 3 = 1.1$
- The privacy-loss expenditure ( $\epsilon$ -differential privacy) is 1.1
- Sources: Warner (1965) [[link](#)] and Greenberg, Abdel-Latif, Simmons, and Horvitz (1969) [[link](#)]. SDL uses: Fienberg and Steele (1998) [[link](#)], Du and Zhan (2003) [[link](#)] and Erlingsson, Vasyl and Korolova (2014) [[link](#)].

# What Happens to Data Quality?

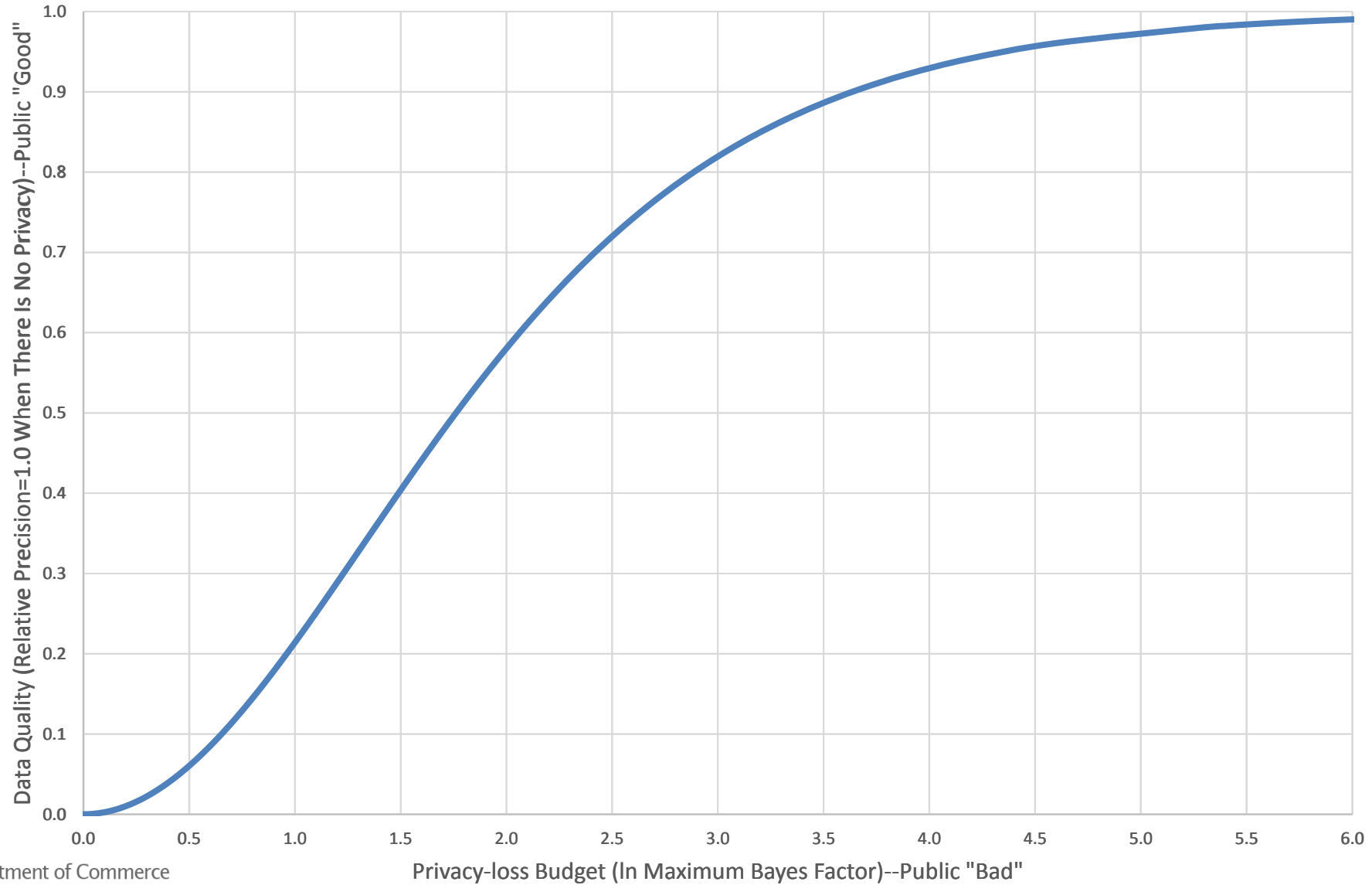
- Use relative sampling precision

$$\text{Rel. Precision} = \frac{\{Pr[\text{Ask Sensitive } Q]\}^2 \frac{n}{\theta(1-\theta)}}{\frac{n}{\theta(1-\theta)}} = \left\{\frac{1}{2}\right\}^2 = 0.25$$

- If
  - Privacy loss is  $\ln 3$
  - Then, relative sampling precision is 25% of the most accurate estimator



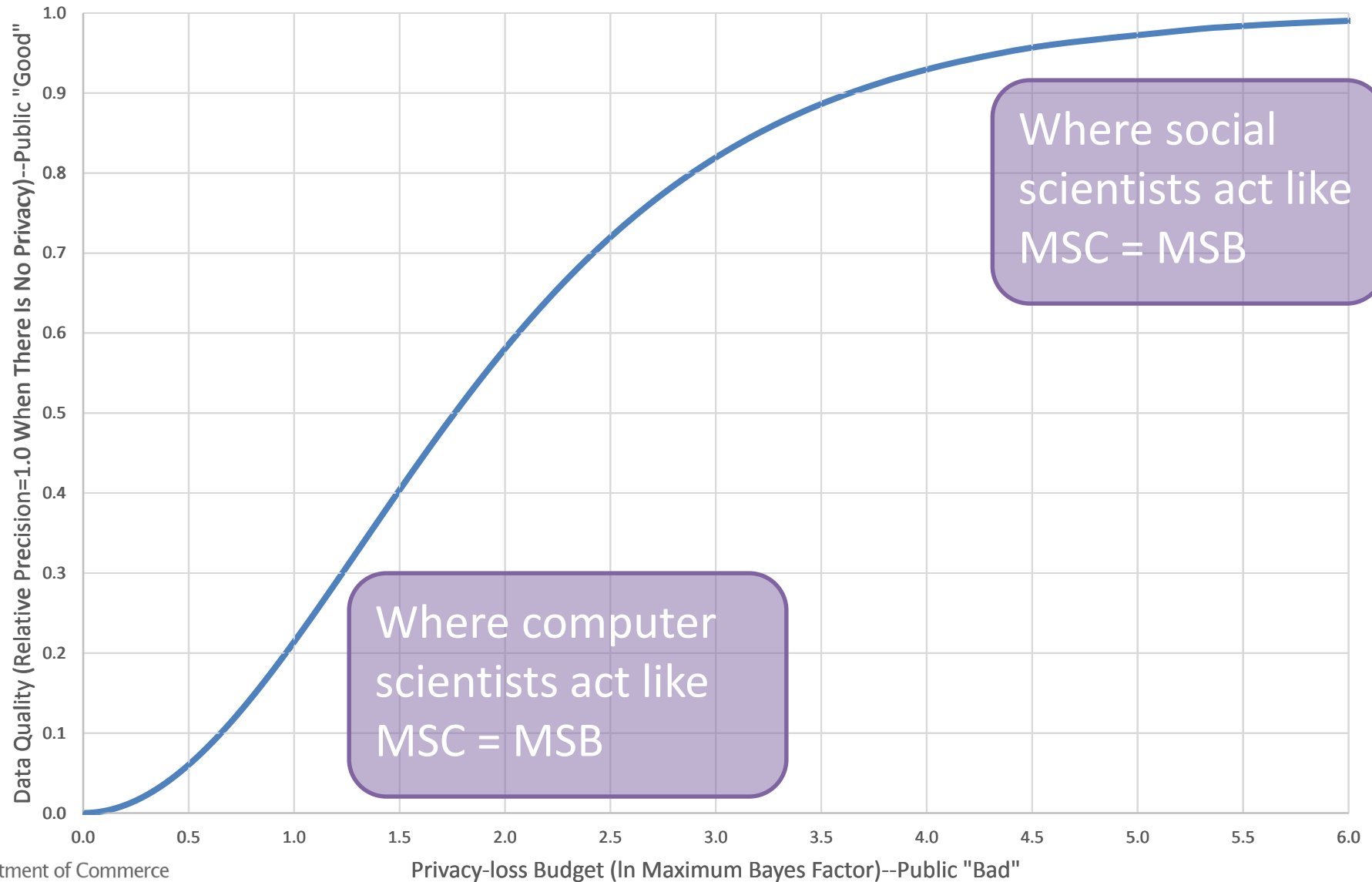
## Production Possibilities Frontier/Risk-Utility/Receiver Operation Characteristics for Statistical Disclosure Limitation via Randomized Response



# Disclosure Limitation is Technology

- The price of increasing data quality (public “good”) in terms of increased privacy loss (public “bad”) is the slope of the technology frontier:
  - Economics: [Production Possibilities Frontier \(Risk-Return in finance\)](#)
  - Forecasting models: [Receiver Operating Characteristics Curve](#)
  - Statistical Disclosure Limitation: [Risk-Utility Curve \(with risk on the x-axis\)](#)
- All exactly the same thing
- None able to select an optimal point

## Production Possibilities Frontier/Risk-Utility/Receiver Operation Characteristics for Statistical Disclosure Limitation via Randomized Response



# Some Examples

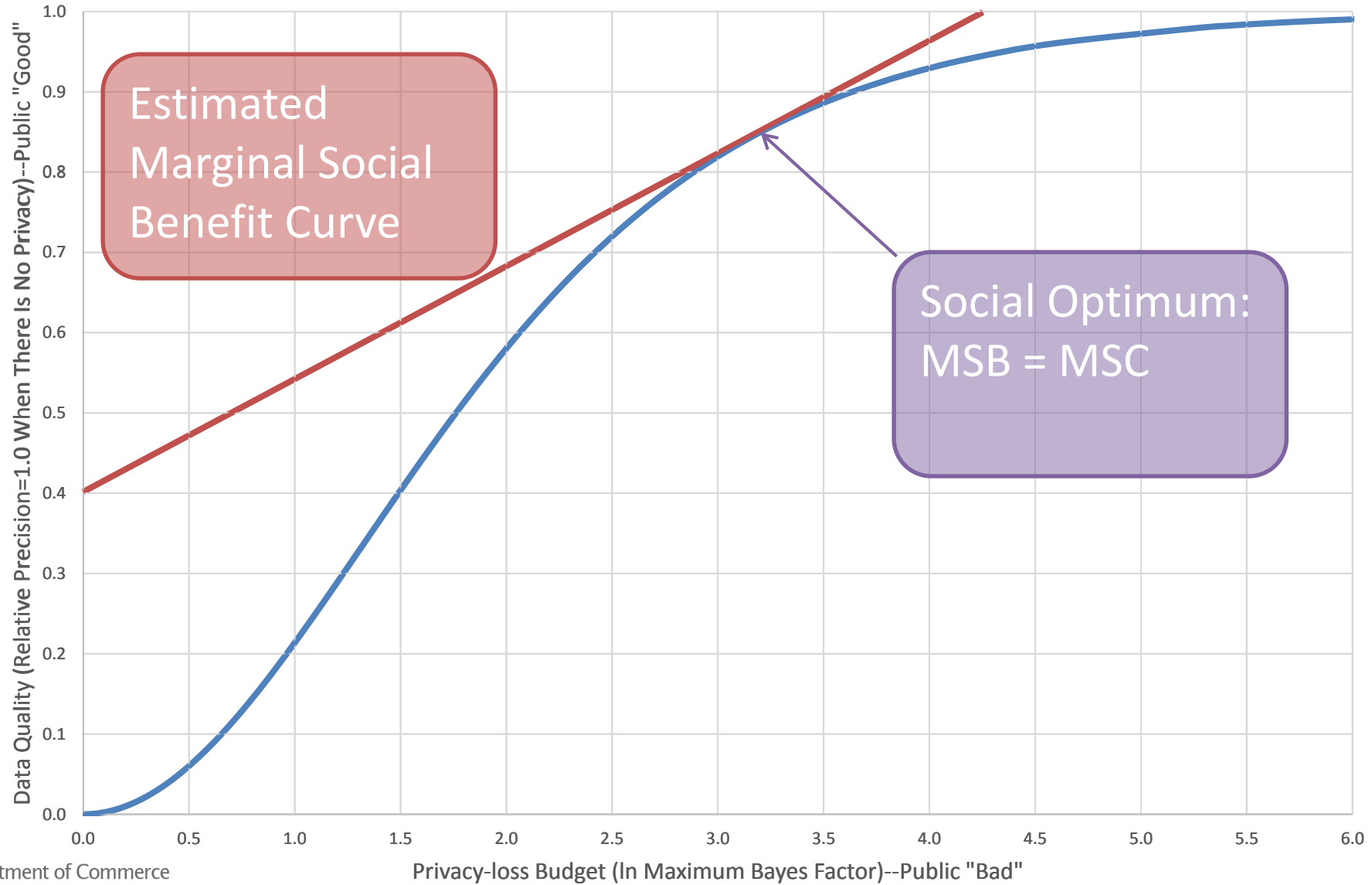
- Dwork (2008): “The parameter  $e$  in Definition 1 is public. The choice of  $e$  is essentially a social question and is beyond the scope of this paper.” [[link](#), p. 3]
- Dwork (2011): “The parameter  $e$  is public, and its selection is a social question. We tend to think of  $e$  as, say, 0.01, 0.1, or in some cases,  $\ln 2$  or  $\ln 3$ .” [[link](#), p. 91]
- In OnTheMap,  $e = 8.9$ , was required to produce tract-level estimates with acceptable accuracy

# How to Think about the Social Choice Problem

- The marginal social benefit is the sum of all citizens' willingness-to-pay for data quality with increased privacy loss
- Can be estimated from survey data
- The next slide shows how

See Abowd and Schmutte (2015) [[link](#)].

# Production Possibilities Frontier/Risk-Utility/Receiver Operation Characteristics for Statistical Disclosure Limitation via Randomized Response



# How to Prove That a Privacy-loss Budget Was Respected

- Must quantify the privacy-loss expenditure of each publication
- The collection of the algorithms taken altogether must satisfy the privacy-loss budget
- Requires methods that compose

# How to Prove That the Algorithms are Resistant to All Future Attacks

- Information environment is changing much faster than before
- *It may no longer be reasonable to assert that a product is empirically safe given best-practice disclosure limitation prior to its release*
- Formal privacy models replace empirical assessment with designed protection
- Resistance to all future attacks is a property of the design



# The Silver Lining

- American Statistical Association on p-values [[link](#)]
- Call for more nuanced use
- Data analysis conducted using privacy-preserving methods:
  - Control the false discovery rate
  - Reduce inferential errors due to multiple comparisons
  - Examples: Erlingsson, Vasyl and Korolova (2014) [[link](#)]; Dwork et al. (2015) [[link](#)]; Apple (2016) [[link](#)]

# A Long Row to Hoe

- Concerted research and engineering effort needed to bring disclosure limitation into the 21<sup>st</sup> century
- Scientific integrity requires that we tackle this challenge
- First step is experimentation with the technologies known to work:
  - Synthetic data with validation using formally private synthesizers
  - Privacy-preserving data analysis via pre-specified query systems

# Thank you.

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